

Appl. No. 10/707,645  
Amdt. dated December 29, 2006  
Reply to Office action of December 14, 2006

**Amendments to the Claims:**

This listing of claims will replace all prior versions and listings of claims in the application:

**Listing of Claims:**

1-13 (cancelled).

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14 (currently amended): A control circuit of memory address decoding for determining whether a given address is located in one of a plurality of sections, each section having at least one memory unit and each memory unit having a unique corresponding address, the corresponding address using the binary system, the control circuit comprising:

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an access module for receiving the given address;

a sorting module for making the corresponding address of the section with greater size smaller than the corresponding address of the section with smaller size, and if the size of a first section is equal to the size of a second section, the first and the second sections are capable of being swapped; and

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a comparing module for building a bit-pattern for each section based on its corresponding addresses and sending a plurality of comparison signals after comparing the given address with those of each bit-pattern, the comparing module comprising a plurality of comparing units, each comparing unit comprising a plurality of first level AND gates, a plurality of NXOR gates, and a second level AND gate, each of the first level AND gates having two inputs for respectively receiving a mask bit generated from the bit-patterns and an associated bit of the given address, each of the NXOR gates having two inputs for respectively receiving the output of one of the first level AND gates and a standard address generated from the bit-patterns, the inputs of the second level  
20 AND gate being connected to the outputs of the NXOR gates and thereby sending out the comparison signals.

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15 (original): The control circuit of claim 14 further comprising a logic module responsible for receiving the comparison signals and sending a decoding result to determine the given address is located in one of the sections.

5 16 (original): The control circuit of claim 14 wherein the sections are a plurality of memory modules.

17 (previously presented): The control circuit of claim 14 wherein the single bit-pattern is built for each section in the comparing module, the bit-pattern consisting of all common  
10 bits of the corresponding addresses in each section.

18 (cancelled).

19 (currently amended): A memory address decoding method for determining an  
15 objective section of a given address in a memory, wherein the memory is formed by at least one section with at least one memory unit, the method comprising:  
assigning an address to each memory unit according to the memory size of the section;  
obtaining at least one bit-pattern of each section according to the common pattern  
20 common-rules of the bit of the addresses; and  
comparing the given address with each bit-pattern to determine the objective section of the given address, and sending a plurality of comparison signals after  
comparing the given address with those of each bit-pattern;  
wherein the addresses of the memory units located in the section with greatest  
25 size are firstly assigned, and the addresses of the memory units located in the section  
with smallest size are lastly assigned.

20 (previously presented): The method of claim 19 wherein the section is formed by at

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least one memory module.

21 (cancelled).

5 22 (previously presented): The method of claim 19 wherein the addresses of the memory units located in the section with greater size are smaller than the addresses of memory units located in the section with smaller size.

10 23 (previously presented): The method of claim 19 wherein the bit-patterns are exclusive to each other.

24 (previously presented): The method of claim 19 wherein the bit-pattern is obtained by all common bits of the address in each section.

15 25 (previously presented): The method of claim 19 wherein the bit-pattern is obtained by partial common bits of the address in each section.

20 26 (previously presented): The method of claim 19 the given address is located in the objective section when certain bits of the given address completely match the bit-pattern of the objective section.

27 (previously presented): The method of claim 19 wherein the size of each section is a power of 2.

25 28 (currently amended): A memory address decoding method for determining an objective section of a given address in a memory, wherein the memory is formed by at least one section with at least one memory unit, the method comprising:  
    sorting the sections according to the memory size of each section;

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assigning an address to each memory unit of each section, wherein the addresses of  
the memory units located in the section with greater size are smaller than the  
addresses of memory units located in the section with smaller size;  
obtaining at least one bit-pattern of each section according to the common pattern  
5 ~~common~~ rules of the bit of the addresses; and  
comparing the given address with each bit-pattern to determine the objective section  
of the given address, and sending a plurality of comparison signals after  
comparing the given address with those of each bit-pattern;  
wherein if the memory size of a first section is ~~substantially~~ equal to a second  
10 section, the addresses of the first section and the second section are swappable.

29 (previously presented): The method of claim 28 wherein the section is formed by at  
least one memory module.

15 30 (previously presented): The method of claim 28 wherein the bit-patterns are exclusive  
to each other.

31 (previously presented): The method of claim 28 wherein the bit-pattern is obtained by  
all common bits of the address in each section.

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32 (previously presented): The method of claim 28 wherein the bit-pattern is obtained by  
partial common bits of the address in each section.

25 33 (previously presented): The method of claim 28 wherein the given address is located in  
the objective section when certain bits of the given address completely match the  
bit-pattern of the objective section.